



Relations

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A New Universe of Discourse

Example 1. English Obstruents

Manner	Location	<i>Lab.</i>	<i>LabVel.</i>	<i>Dent.</i>	<i>Alv.</i>	<i>Alv.-Pal.</i>	<i>Vel.</i>	<i>Glott.</i>
Stops	Oral	<i>p/b</i>			<i>t/d</i>		<i>k/g</i>	
	Nasal	<i>m</i>			<i>n</i>		<i>ŋ</i>	
Fricatives	Oral		<i>f/v</i>	<i>θ/ð</i>	<i>s/z</i>	<i>ʃ/ʒ</i>		<i>h</i>

StopFric | SameVoice | SamePlace |

New Sets

Obstruents = { p, b, t, d, k, g, m, n, η, s, z, θ, ð, ʃ, ʒ, h } |Obstruents| = 16

Stop = {p, b, t, d, k, g, m, n, η} |Stops| = 9

Fric = {θ, ð, s, z, ʃ, ʒ, h} |Fricatives| = 7

Nas = {m, n, η} |Nasals| = 3

Oral = {p, b, t, d, s, z, k, g, θ, ð, ʃ, ʒ, h} |Orals| = 13

OralStop = {p, b, t, d, k, g} |OralStops| = 6

Introducing Ordered Pairs

- Ordered Pairs

$$\langle a, b \rangle$$

- Order matters

$$\langle a, b \rangle \neq \langle b, a \rangle$$

- Not like sets

$$\{a, b\} = \{b, a\}$$

Ordered pairs versus sets: More differences

- The same object can occur more than once

$$\langle a, a \rangle \neq a \neq \langle a \rangle \neq \{a\}$$

- Unlike sets

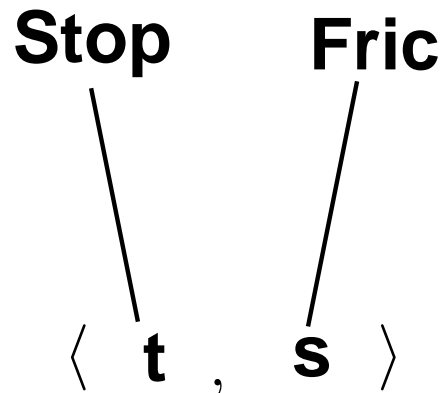
$$\{a\} = \{a, a\} = \{a, a, a\} \dots$$

Sets of Ordered pairs [Relations]

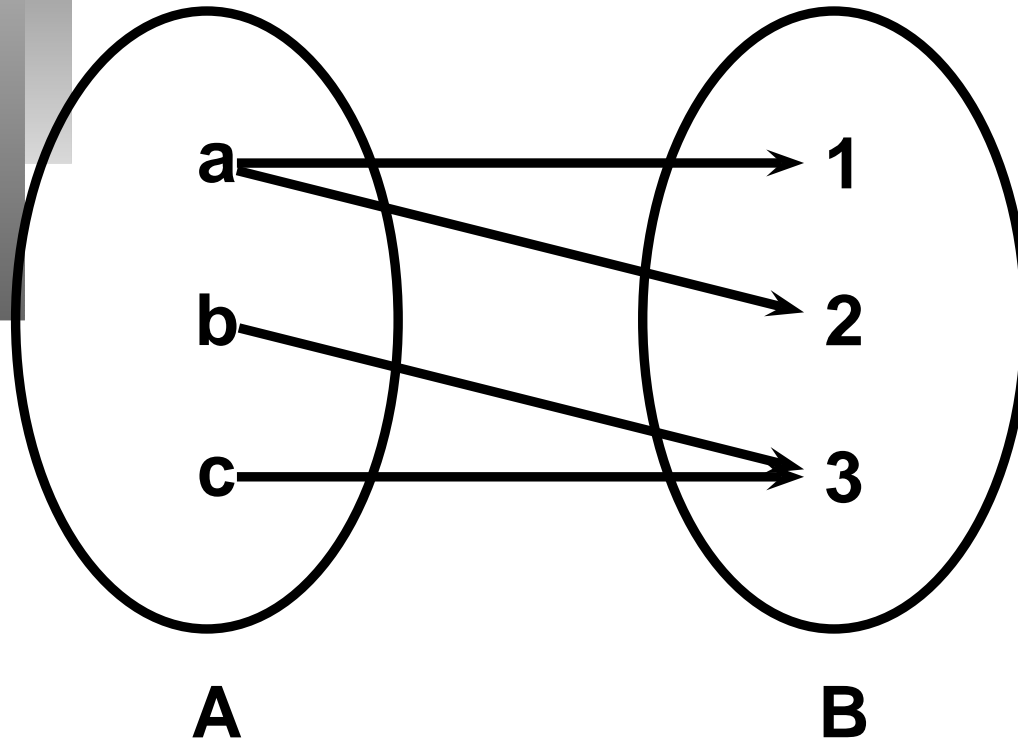
- A set of ordered pairs

$$\text{StopFric} = \{ \langle t, s \rangle, \langle t, z \rangle, \langle d, s \rangle, \\ \langle d, z \rangle, \langle n, s \rangle, \langle n, z \rangle \}$$

- Each *first* member is from Stop, each *second* member is from Fric:



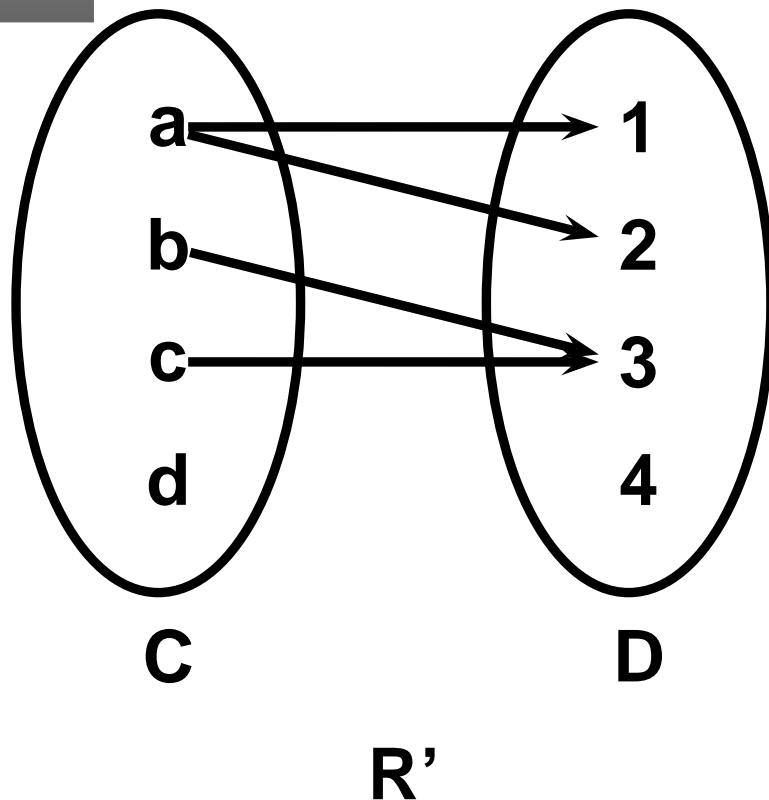
A Relation named R from A to B



$$R = \{\langle a, 1 \rangle, \langle a, 2 \rangle, \langle b, 3 \rangle, \langle c, 3 \rangle\}$$

Domain and Range of a relation

A relation R' from C to D The subsets of C and D actually standing in R' are called the domain and range of R' .



$$\text{Dom}(R') = \{a, b, c\}$$

$$\text{Range}(R') = \{1, 2, 3\}$$

Cartesian Products

$\text{Stop} \times \text{Fric} = \{ \langle x, y \rangle \mid x \in \text{Stop and } y \in \text{Fric} \}$

$\{ \langle p, \theta \rangle, \langle p, \eth \rangle, \langle p, s \rangle, \langle p, z \rangle, \langle p, \int \rangle, \langle p, \mathfrak{Z} \rangle, \langle p, x \rangle,$
 $\langle t, \theta \rangle, \langle t, \eth \rangle, \langle t, s \rangle, \langle t, z \rangle, \langle t, \int \rangle, \langle t, \mathfrak{Z} \rangle, \langle t, x \rangle,$
 $\langle k, \theta \rangle, \langle k, \eth \rangle, \langle k, s \rangle, \langle k, z \rangle, \langle k, \int \rangle, \langle k, \mathfrak{Z} \rangle, \langle k, x \rangle,$
 $\langle b, \theta \rangle, \langle b, \eth \rangle, \langle b, s \rangle, \langle b, z \rangle, \langle b, \int \rangle, \langle b, \mathfrak{Z} \rangle, \langle b, x \rangle,$
 $\langle d, \theta \rangle, \langle d, \eth \rangle, \langle d, s \rangle, \langle d, z \rangle, \langle d, \int \rangle, \langle d, \mathfrak{Z} \rangle, \langle d, x \rangle,$
 $\langle g, \theta \rangle, \langle g, \eth \rangle, \langle g, s \rangle, \langle g, z \rangle, \langle g, \int \rangle, \langle g, \mathfrak{Z} \rangle, \langle g, x \rangle,$
 $\langle m, \theta \rangle, \langle m, \eth \rangle, \langle m, s \rangle, \langle m, z \rangle, \langle m, \int \rangle, \langle m, \mathfrak{Z} \rangle, \langle m, x \rangle,$
 $\langle n, \theta \rangle, \langle n, \eth \rangle, \langle n, s \rangle, \langle n, z \rangle, \langle n, \int \rangle, \langle n, \mathfrak{Z} \rangle, \langle n, x \rangle,$
 $\langle \eta, \theta \rangle, \langle \eta, \eth \rangle, \langle \eta, s \rangle, \langle \eta, z \rangle, \langle \eta, \int \rangle, \langle \eta, \mathfrak{Z} \rangle, \langle \eta, x \rangle \}$

Order in Cartesian Products

- Stop \times Fric is order dependent, because the ordered-pairs in the set are:

$$\text{Fric} \times \text{Stop} \neq \text{Stop} \times \text{Fric}$$

$$\langle s, t \rangle \in \text{Fric} \times \text{Stop}$$

$$\langle s, t \rangle \notin \text{Stop} \times \text{Fric}$$

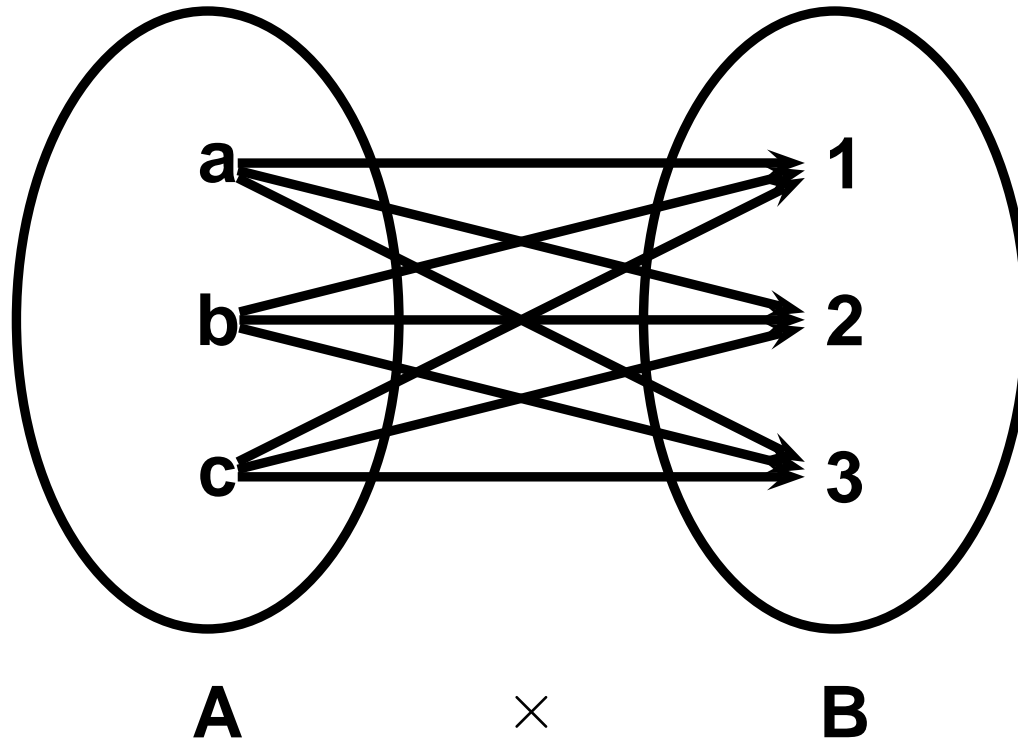
$$\langle t, s \rangle \notin \text{Fric} \times \text{Stop}$$

$$\langle t, s \rangle \in \text{Stop} \times \text{Fric}$$

- Stop \times Fric is a special case of a relation from Stop to Fric. For any such relation R,

$$R \subseteq \text{Stop} \times \text{Fric}$$

The Cartesian product of A and B



$$A \times B = \{ \langle a, 1 \rangle, \langle a, 2 \rangle, \langle a, 3 \rangle, \\ \langle b, 1 \rangle, \langle b, 2 \rangle, \langle b, 3 \rangle, \\ \langle c, 1 \rangle, \langle c, 2 \rangle, \langle c, 3 \rangle \}$$

Cardinality of Cartesian Product

Observation: To construct $A \times B$ we need to pair each member of A with all $|B|$ members of B . Since we have to do that $|A|$ times

$$|A \times B| = |A| \times |B|$$

Identity Relation

A special relation from any set A to itself is the identity relation:

$$I_A = \{\langle x, x \rangle \mid x \in A\}$$

SamePlace, SameVoicing

Relations can be defined by listing OR by predicates, just as sets can be. Reassuring because they are sets.

Example 2. *SamePlace*

$$\text{SamePlace} = \{ \langle x, y \rangle \mid \langle x, y \rangle \in \text{Obstruents} \times \text{Obstruents} \text{ and } x \text{ and } y \text{ are articulated in the same place.} \}$$

Example 3. *SameVoicing*

$$\text{SameVoicing} = \{ \langle x, y \rangle \mid \langle x, y \rangle \in \text{Obstruents} \times \text{Obstruents} \text{ and } x \text{ and } y \text{ have the same voicing.} \}$$

These are relations from Obstruents to Obstruents.

Example 4. *StopFric*

$$\text{StopFric} = \{ \langle x, y \rangle \mid \langle x, y \rangle \in \text{Stop} \times \text{Fric} \text{ and} \\ \text{SamePlace}(x, y) \}$$

Consulting *example 1*, we find :

$$\text{StopFric} = \{ \langle t, s \rangle, \langle t, z \rangle, \langle d, s \rangle, \\ \langle d, z \rangle, \langle n, s \rangle, \langle n, z \rangle \}$$

Definition 1. *Inverse of a Relation*

We define the inverse \mathbf{R}^{-1} of a relation \mathbf{R} as follows:

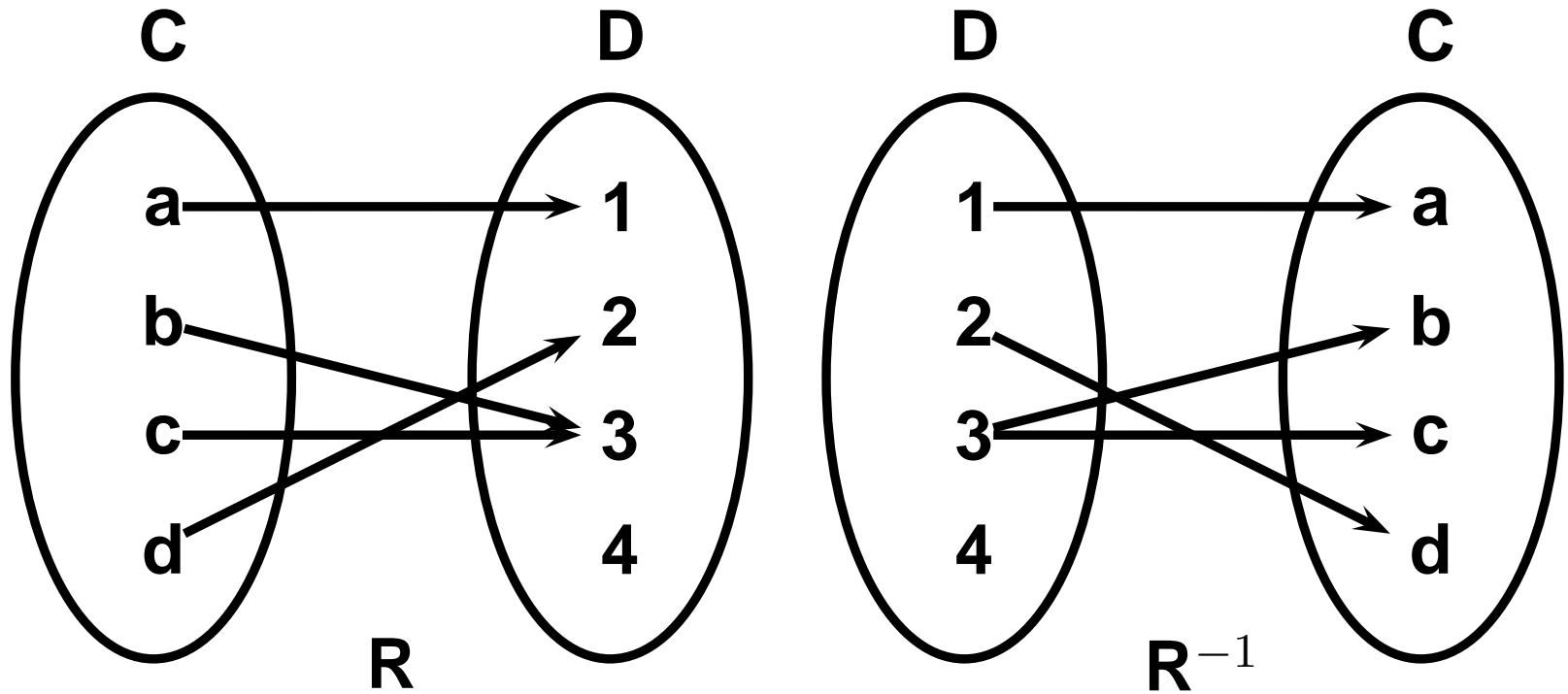
$$\mathbf{R}^{-1} = \{ \langle x, y \rangle \mid \langle y, x \rangle \in \mathbf{R} \}$$

Note that it is frequently the case that:

$$\mathbf{R}^{-1} \cap \mathbf{R} = \emptyset$$

When is it not?

A relation and its inverse



Example 5. Inverse of StopFric

$$\begin{aligned} \text{StopFric}^{-1} &= \{\langle s, t \rangle, \langle z, t \rangle, \langle s, d \rangle, \langle z, d \rangle, \langle s, n \rangle, \langle z, n \rangle\} \\ &= \text{FricStop} \end{aligned}$$

Relation Composition

Definition 2.

$$R_1 \circ R_2 = \{ \langle x, z \rangle \mid \exists y [\langle x, y \rangle \in R_2 \text{ and } \langle y, z \rangle \in R_1] \}$$

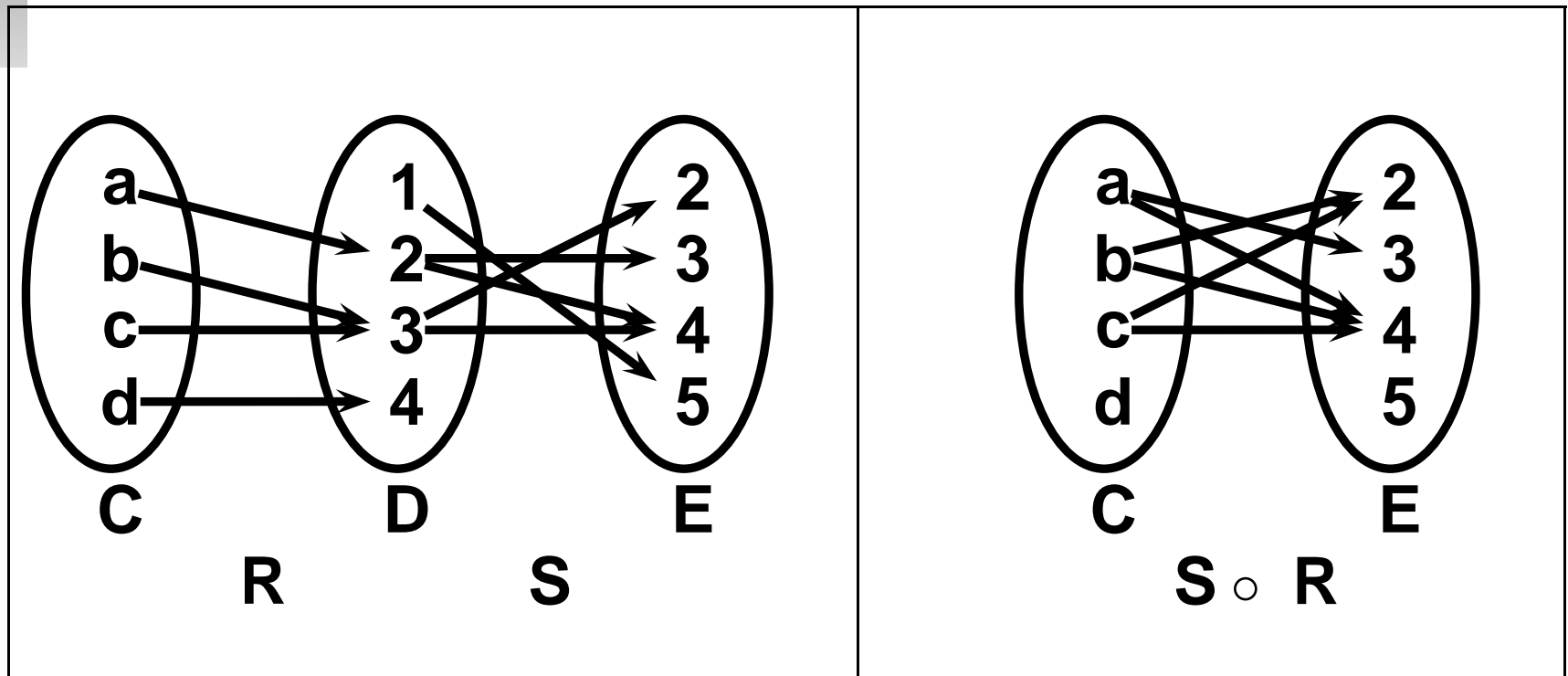
We introduce a new piece of notation here \exists , which means “there exists”.

$$R_1 \circ R_2 = \{ \langle x, z \rangle \mid \text{There exists a } y \text{ such that } \langle x, y \rangle \in R_2 \text{ and } \langle y, z \rangle \in R_1 . \}$$

Note:

If R_2 is from A to B and R_1 is from C to D then $R_1 \circ R_2$ is from A to D .

Composition of relations R and S .



Although $S \circ R$ must be from the same set as R and to the same set as S it does not have to have the same domain as R or the same range as S .

Defining SamePlace using relation composition

Example 6. Places

$Places = \{bilabial, interdental, alveolar, palatal, velar\}$

Example 7. SoundPlace

$SoundPlace = \{\langle x, y \rangle \mid x \in Obstruents \text{ and } y \in Places \text{ and } x \text{ is pronounced at place } y\}$

Example 8. SamePlace defined by composition

$SamePlace = SoundPlace^{-1} \circ SoundPlace$

Example 9. *Following the relation links for one pair*

$$\langle t, \text{alveolar} \rangle \in \text{SoundPlace}$$

Now $\text{SoundPlace}^{-1} \circ \text{SoundPlace}$ includes all pairs $\langle t, x \rangle$ such that

$$\langle \text{alveolar}, x \rangle \in \text{SoundPlace}^{-1}$$

That is:

$$\{\langle \text{alveolar}, t \rangle, \langle \text{alveolar}, d \rangle, \langle \text{alveolar}, s \rangle, \langle \text{alveolar}, z \rangle\}$$

Thus $\text{SoundPlace}^{-1} \circ \text{SoundPlace}$ includes

$$\{\langle t, t \rangle, \langle t, d \rangle, \langle t, s \rangle, \langle t, z \rangle\}$$



Functions

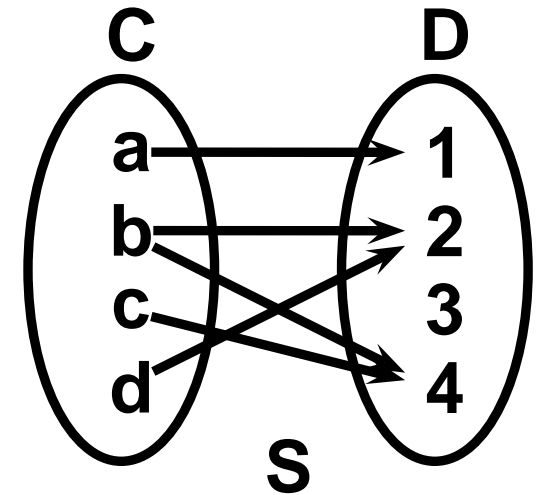
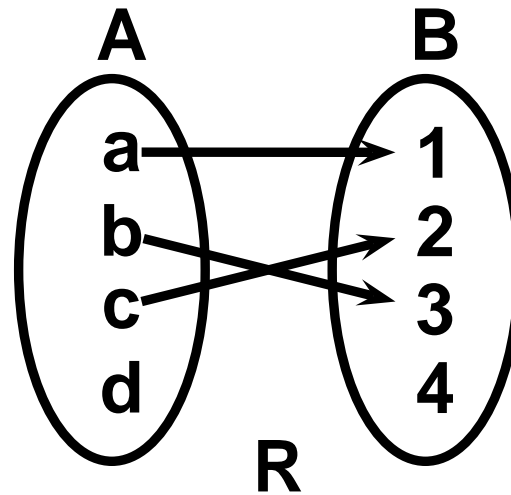
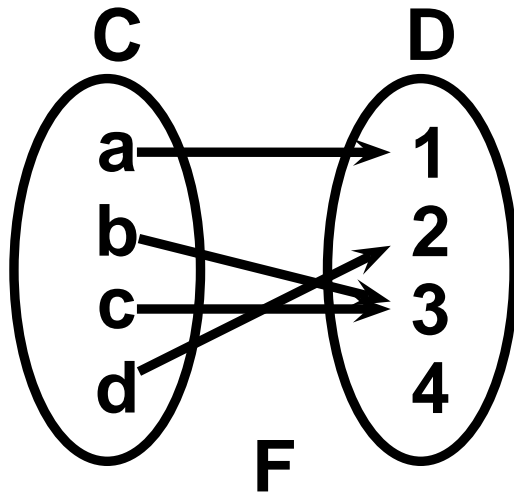
A special case of relations

Definition 3. *Functions*

A relation R from A to B is a function from A to B if and only if

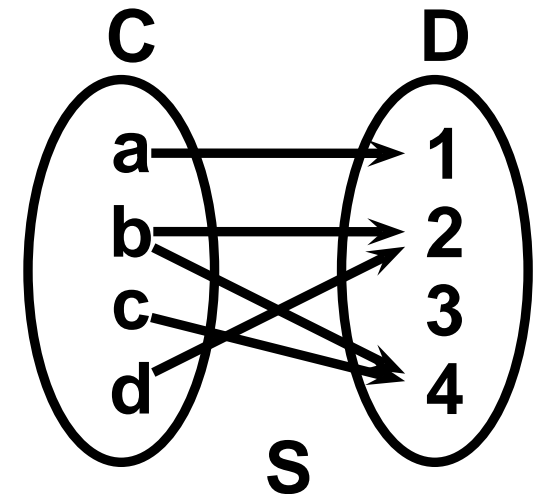
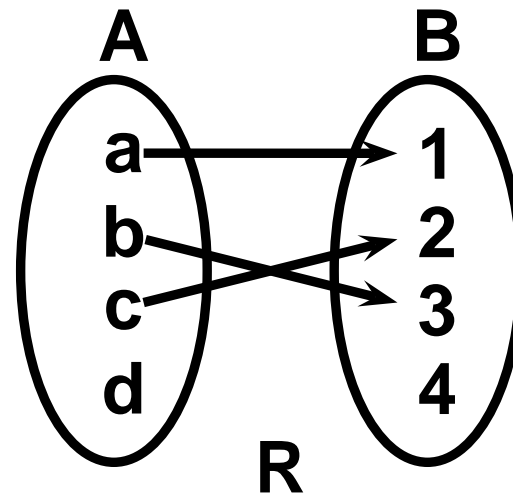
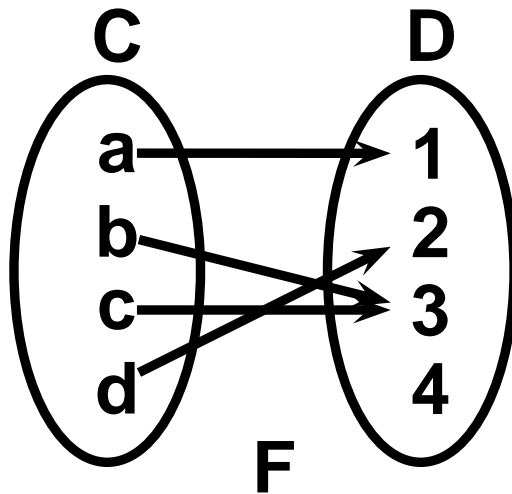
- 1. For each member x of the domain of R , there is a unique member of the range standing in relation R to x .*
- 2. $\text{Dom}(R) = A$.*

Functions and non functions



F
R
S

Functions and non functions

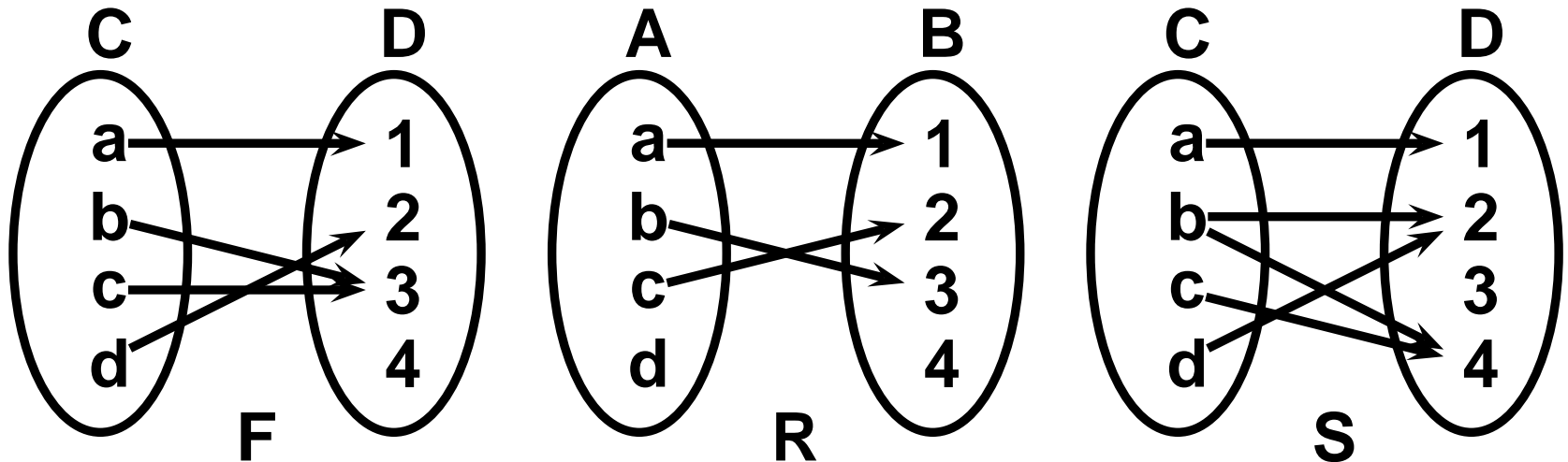


F is a function from C to D.

R

S

Functions and non functions

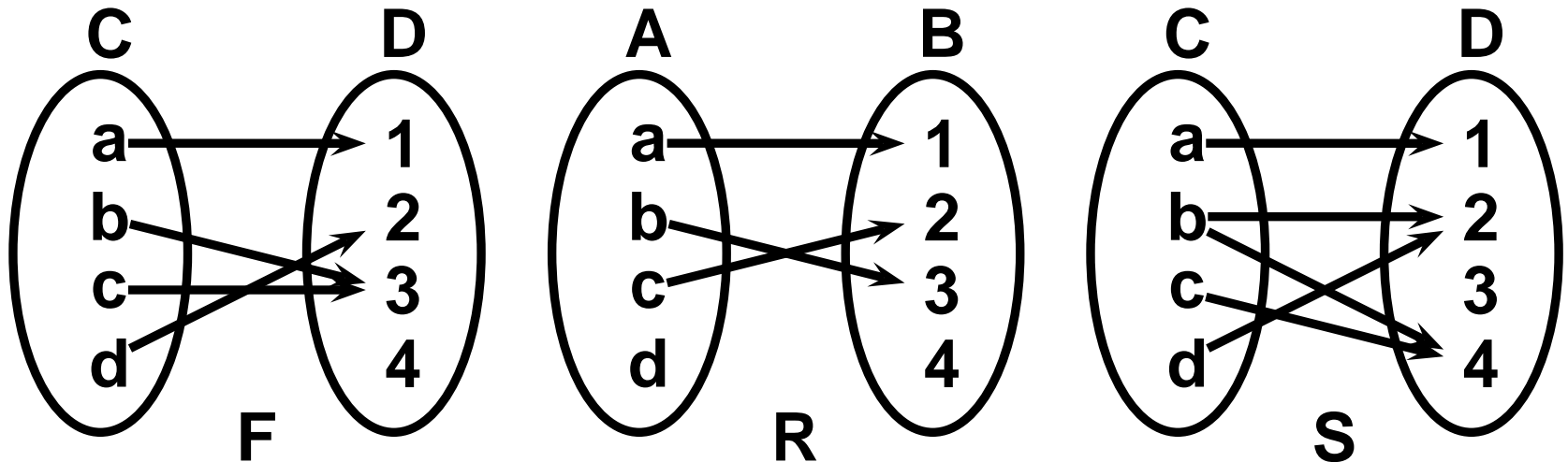


F is a function from C to D.

R is not a function from A to B (Clause 2)

S

Functions and non functions



F is a function from C to D.

R is not a function from A to B (Clause 2)

S is not a function (Clause 1)

Another function

- **SoundPlace** is a function from obstruents to places because each obstruent has exactly one place it is pronounced at.
- **SoundPlace**⁻¹ is NOT a function because each position may have many sounds pronounced at it. For example

$$\langle \textit{alveolar}, s \rangle \in \text{SoundPlace}^{-1}$$

$$\langle \textit{alveolar}, t \rangle \in \text{SoundPlace}^{-1}$$

$$\langle \textit{alveolar}, d \rangle \in \text{SoundPlace}^{-1}$$

$$\langle \textit{alveolar}, z \rangle \in \text{SoundPlace}^{-1}$$

$$\langle \textit{alveolar}, n \rangle \in \text{SoundPlace}^{-1}$$

The successor relation

- The successor relation in arithmetic is a function from integers to integers:

$$s(0) = 1 \quad (1)$$

$$s(1) = 2 \quad (2)$$

$$s(2) = 3 \quad (3)$$

$$\vdots \quad (4)$$

$$s(n) = n + 1 \quad (5)$$

- For each integer there is a unique next number. Every integer HAS a successor and every integer has exactly one.

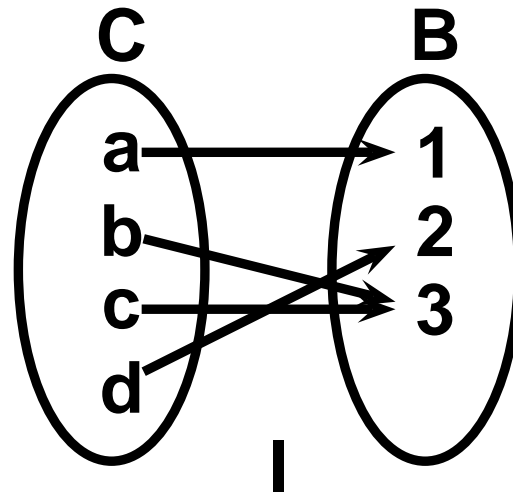
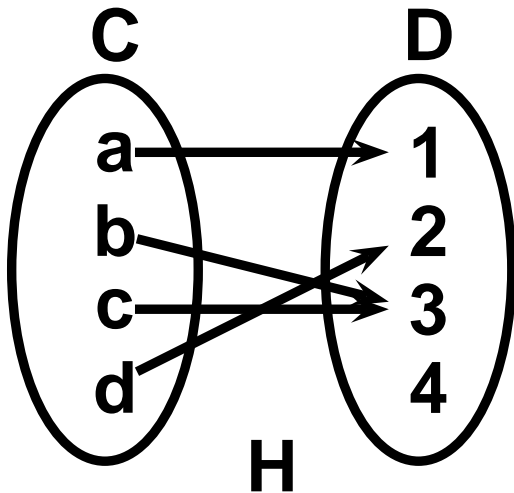
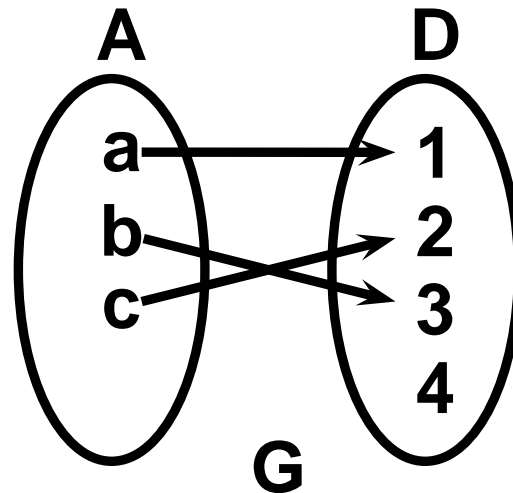
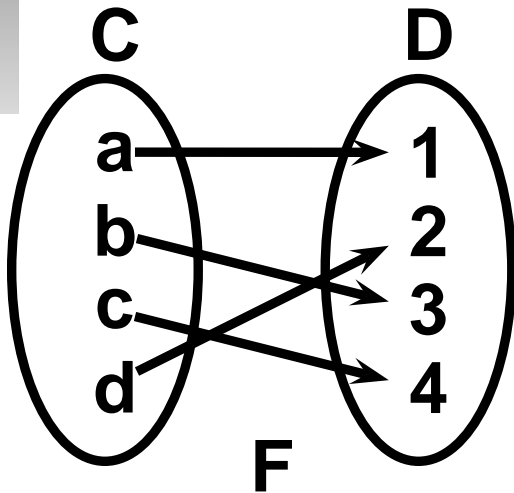
One-to-one functions

Definition 4. *We call a function F from A to B **one-to-one** (or **injective**) if and only if each member of F is paired with exactly one member of B .*

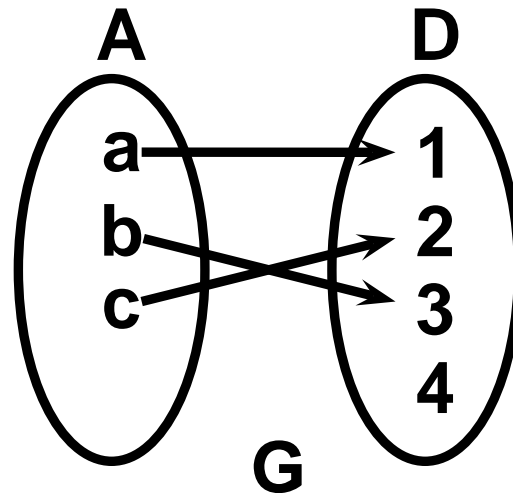
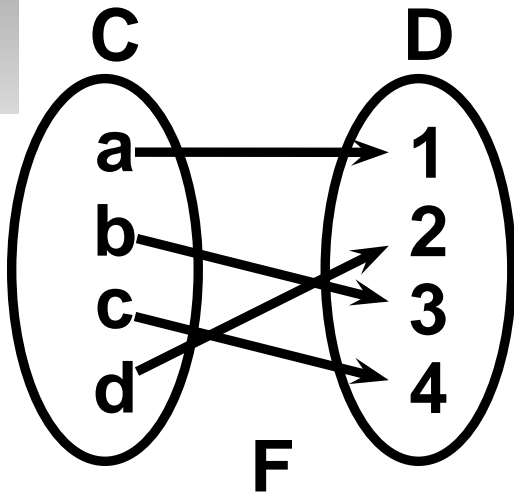
Onto functions

Definition 5. *We call a function **onto** (or surjective) if the range is equal to B . The term into does not mean the opposite of one-to-one. It's just a synonym for to. A function from A to B is also called a function from A into B , with no implied claim about whether it is onto or not.*

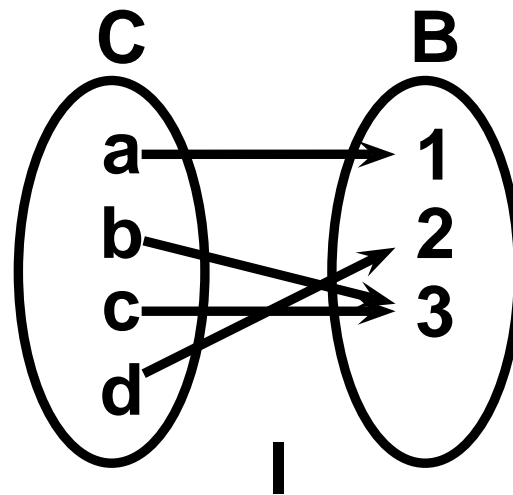
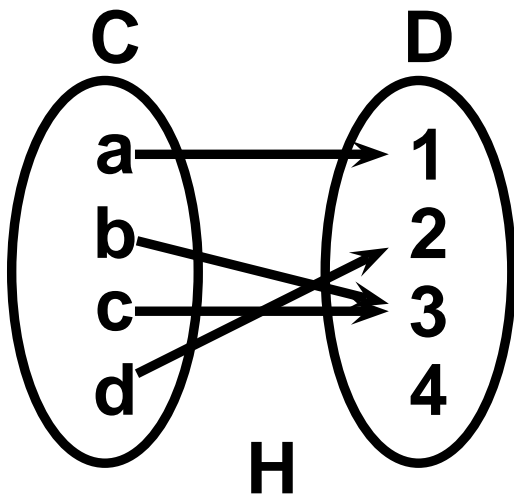
One-to-one and onto functions



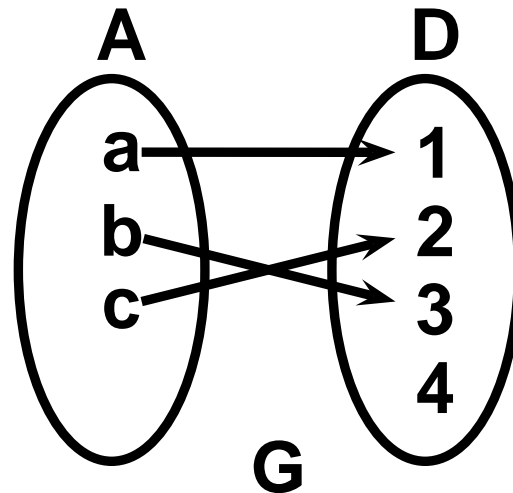
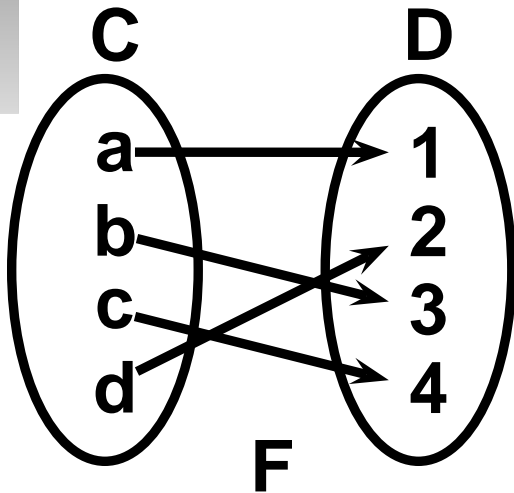
One-to-one and onto functions



F: 1-1, onto

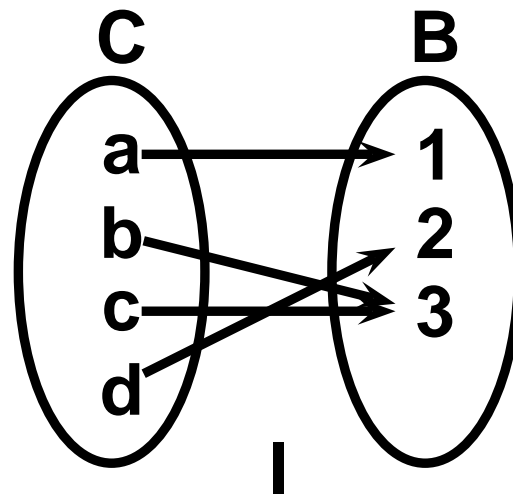
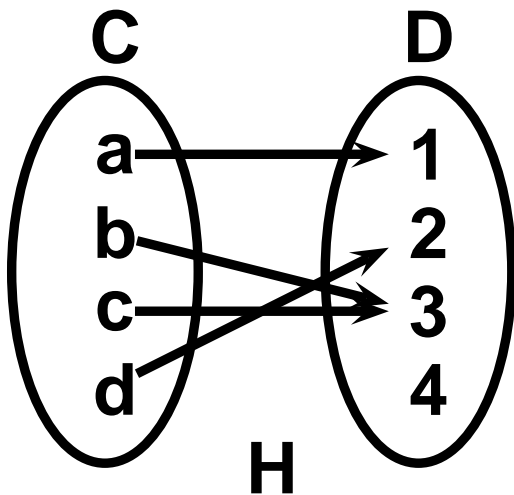


One-to-one and onto functions

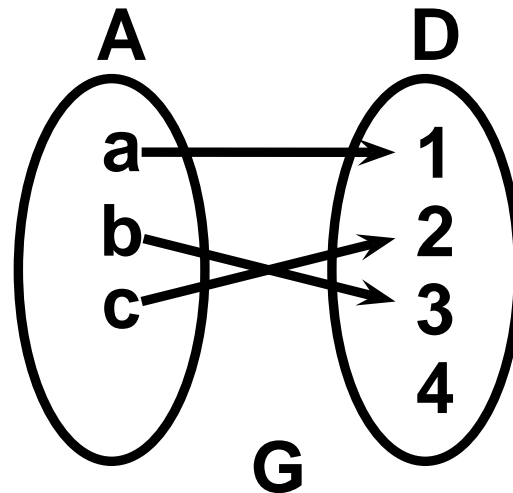
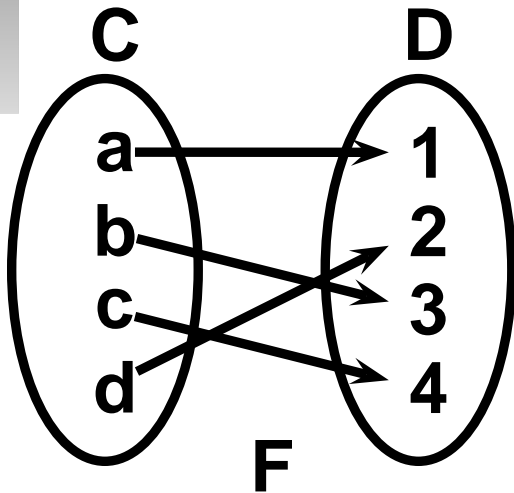


F: 1-1, onto

G: 1-1

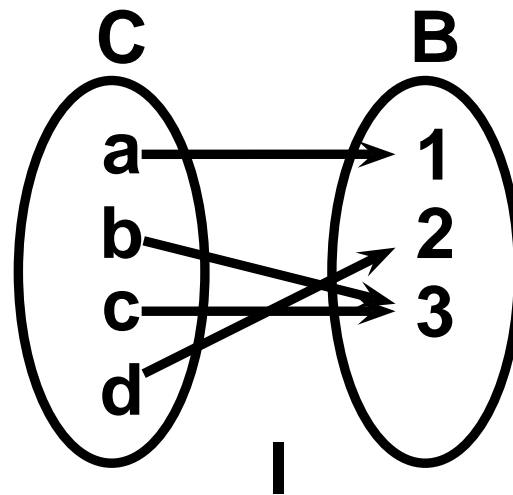
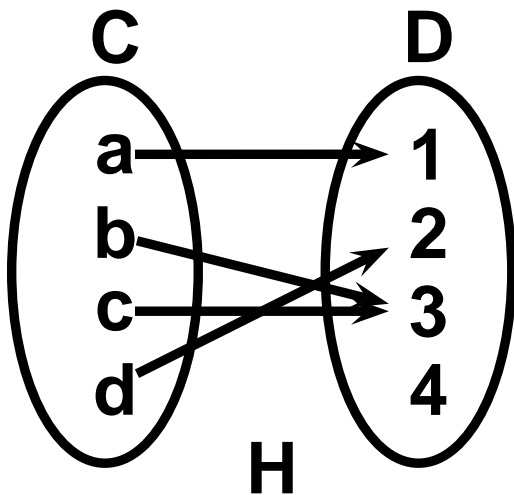


One-to-one and onto functions



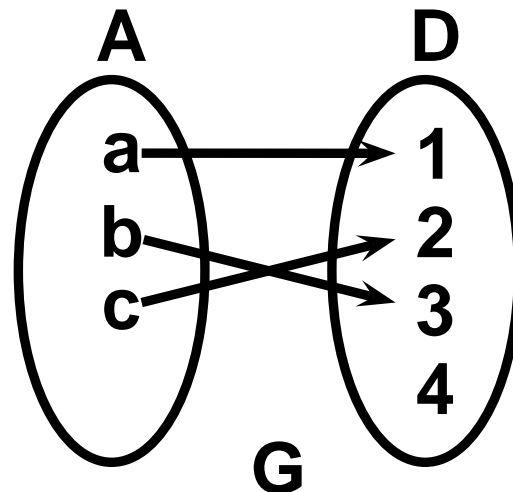
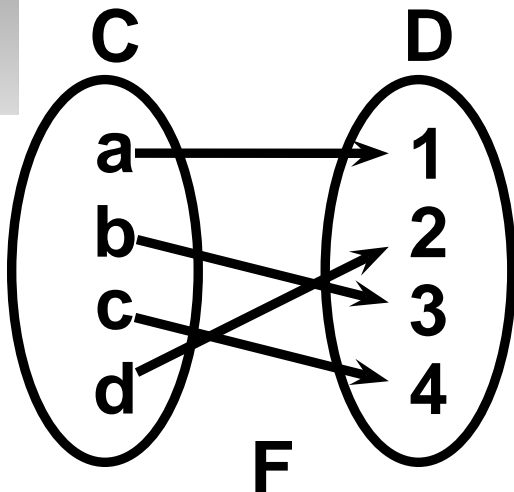
F: 1-1, onto

G: 1-1



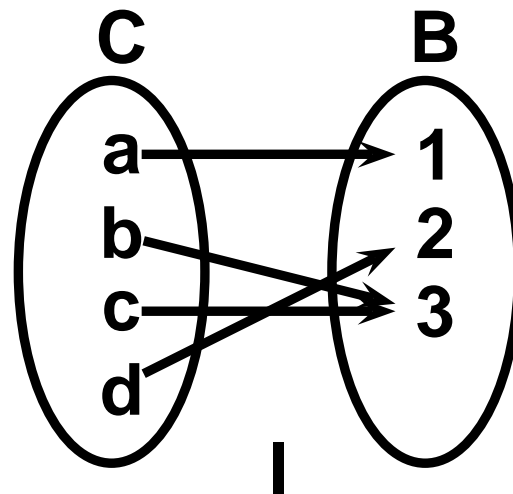
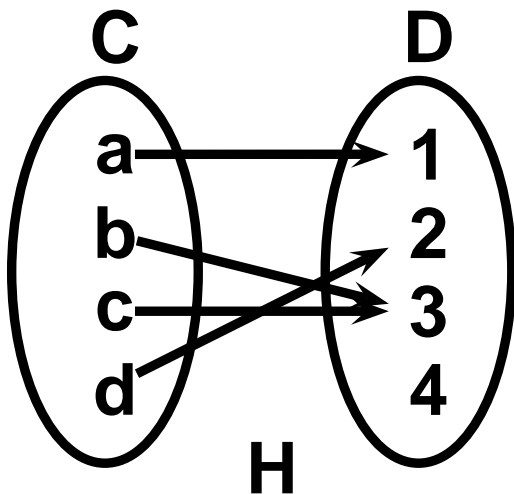
H: neither

One-to-one and onto functions



F: 1-1, onto

G: 1-1



H: neither

I: onto



Some properties of Relations

- We call a pair of the form

$$\langle x, x \rangle$$

a **reflexive** pair.

- Note that some relations have reflexive pairs:

$$\langle t, t \rangle \in \text{same-place}$$

- Some do not.

$$\langle t, t \rangle \notin \text{SoundPlace}$$

Definition 6. *Reflexivity of a relation*
A relation R is reflexive if and only if

$$\forall x \in \text{Dom}(R) \langle x, x \rangle \in R$$

This introduces the symbol \forall (upside-down A). $\forall x$ should be read “for all x ”.

for all x such that $x \in \text{Dom}(R) \langle x, x \rangle \in R$

$\forall x$ is always followed by some statement of a condition that all x 's satisfy. Usually we restrict $\forall x$ to particular set S of x 's: $\forall x \in S$.

Examples of reflexivity

Example 10. *Reflexivity of SamePlace*

The SamePlace relation is reflexive. Every sound is pronounced in the same place as itself.

Example 11. *NonReflexivity of SoundPlace*

The SoundPlace relation is not reflexive. No sound has itself as its place of articulation.

Definition 7. *Irreflexivity of a relation*
A relation R is irreflexive if and only if

$$\forall x \in \text{Dom}(R) \langle x, x \rangle \notin R$$

Example of irreflexivity

Example 12. *Irreflexivity of SoundPlace*

The SoundPlace relation is irreflexive. No sound has itself as its place of articulation. Observation: Any relation from sets A to set B is trivially irreflexive when A and B are disjoint.

Neither Reflexive nor Irreflexive

Example 13. *sum-of-factors (perfect numbers)*
Consider the function *sum-factors*.

$$\text{sum-factors}(x) = y$$

if and only if the sum of all the prime numbers that evenly divide x , excluding x equals y .

sum-factors(2) = 1 1 is the only prime factor of 2

sum-factors(4) = 3 1 and 2 are 4's prime factors and $1+2=3$

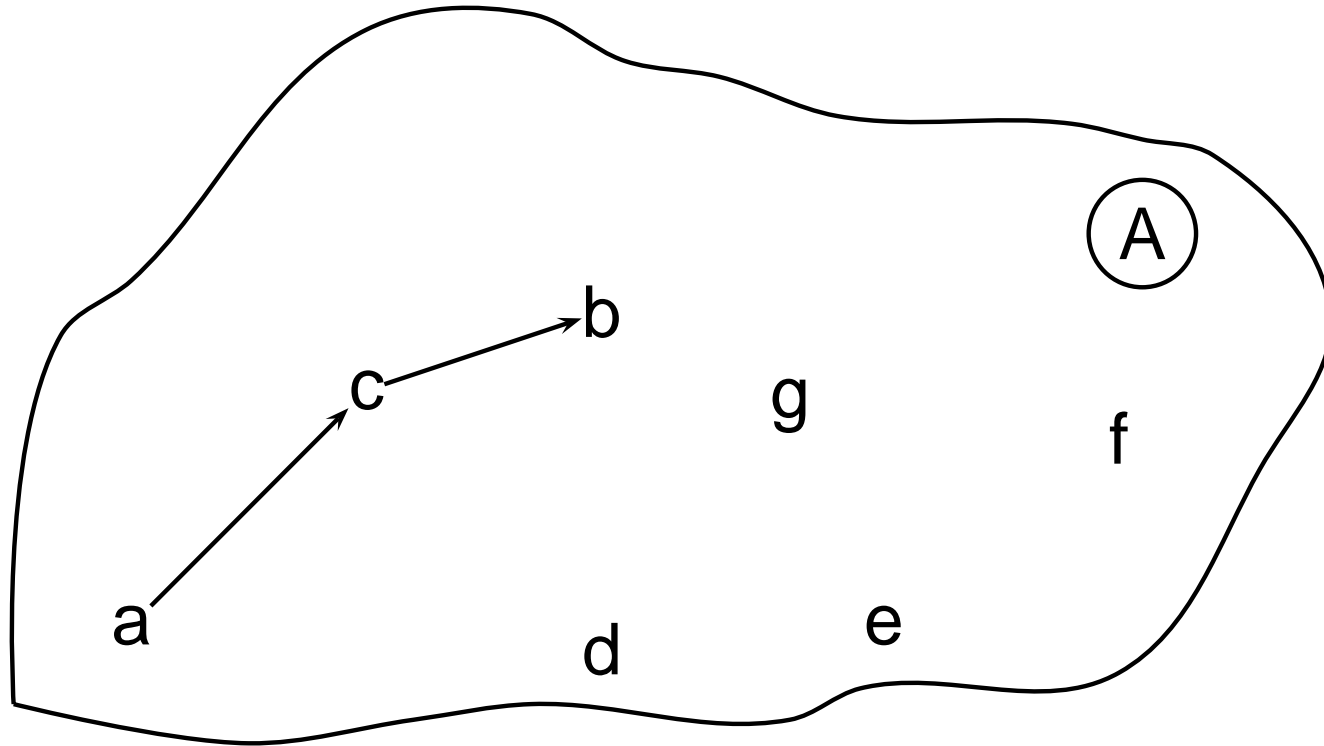
Sum-factors is not reflexive. Is it irreflexive? No, because the prime factors of 6 are 1, 2 and 3 and

$$\text{sum-factors}(6) = 6 \quad 6 = 1 + 2 + 3$$

Definition 8. *Transitivity of a relation*
A relation R is transitive if and only if

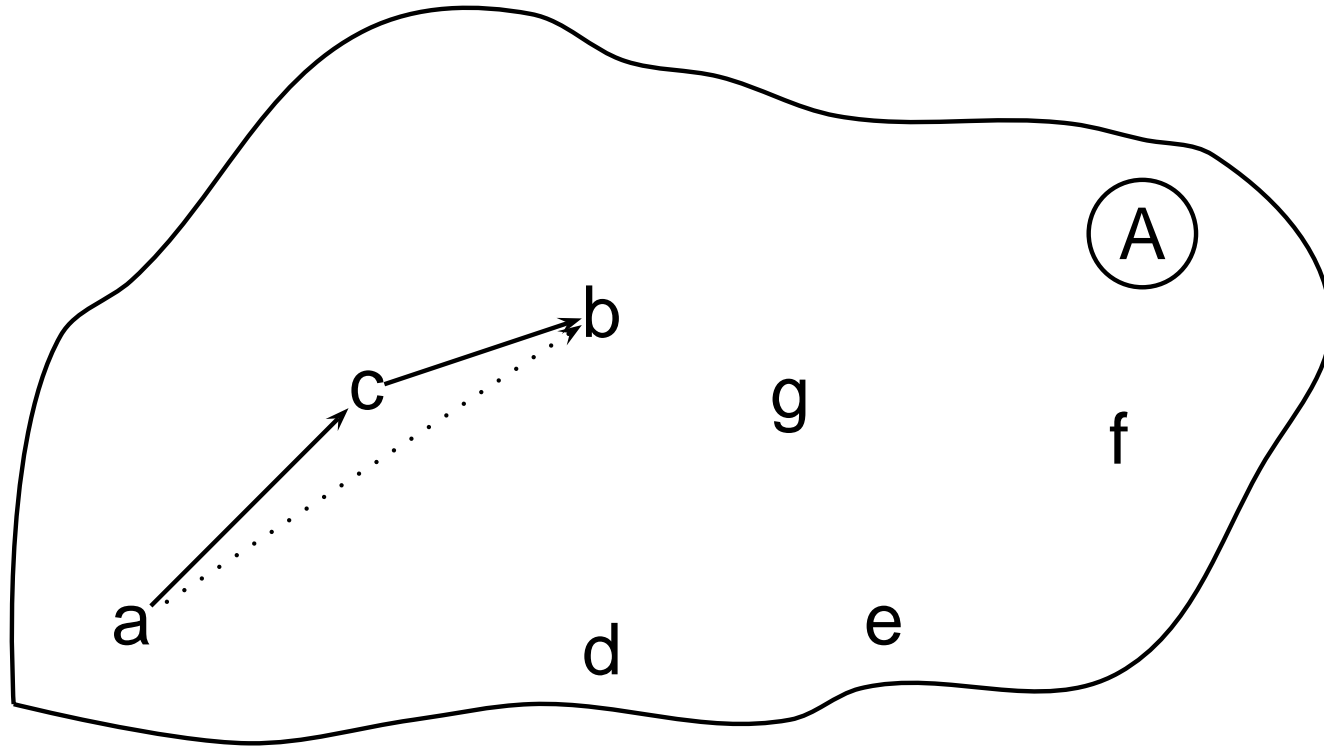
$$\forall \langle x, y \rangle \in R \text{ If } \exists z \langle y, z \rangle \in R \text{ then } \langle x, z \rangle \in R$$

Transitivity: The intuition



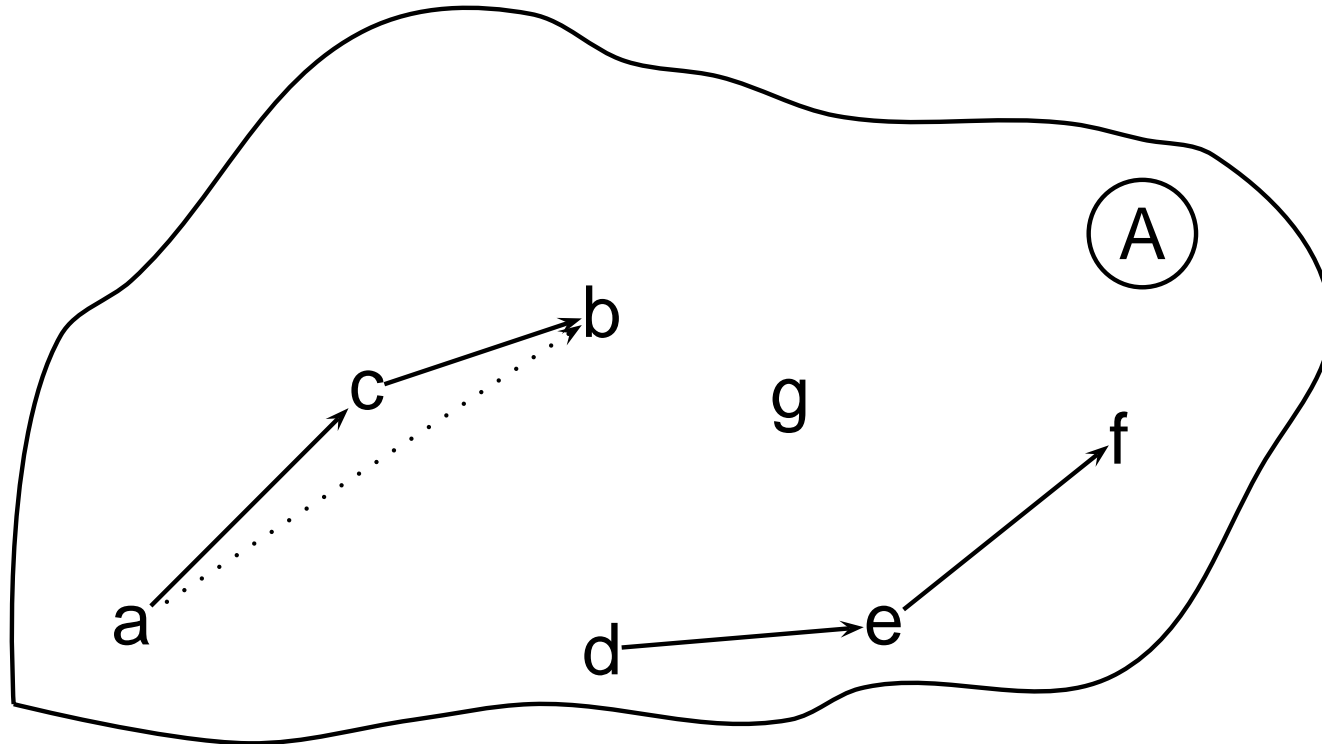
Solid links give a partial picture of a relation R from A to A . If R is transitive, then when the solid links exist, the dashed links must also exist.

Transitivity: The intuition



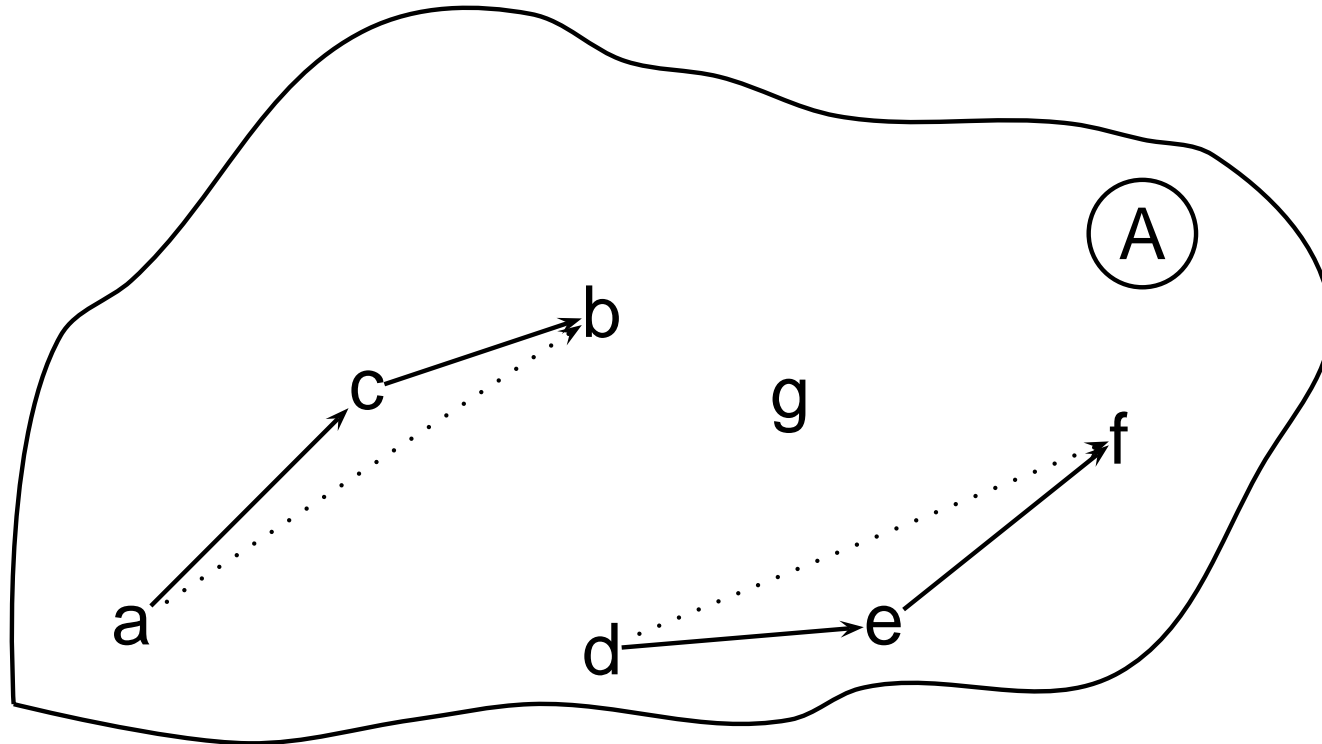
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Transitivity: The intuition



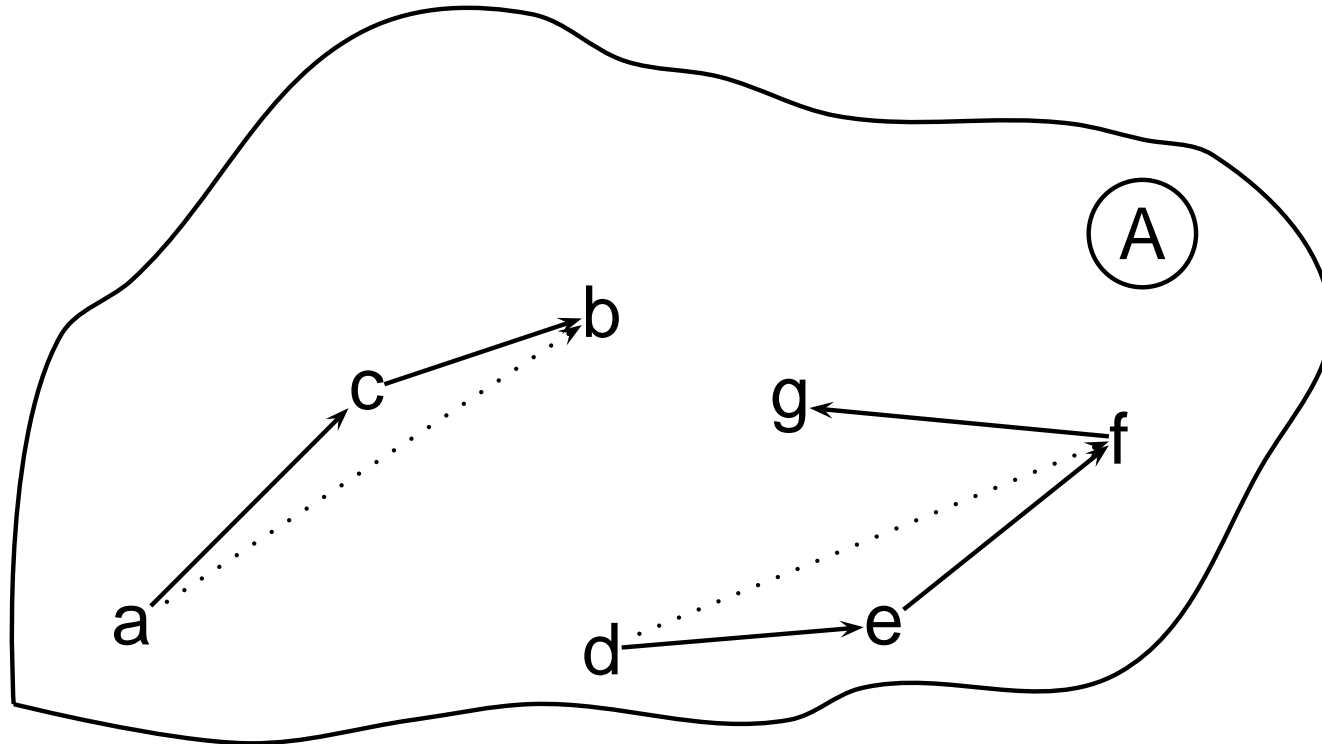
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Transitivity: The intuition



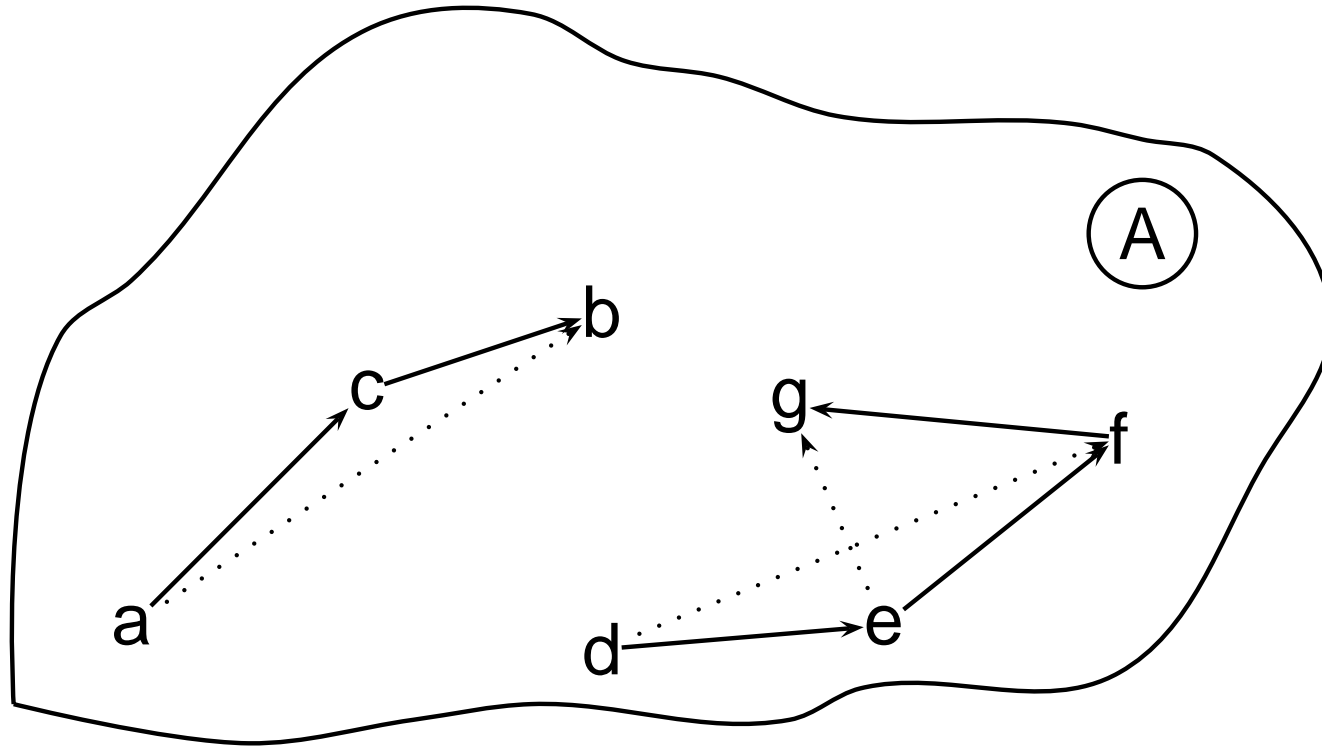
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Transitivity: The intuition



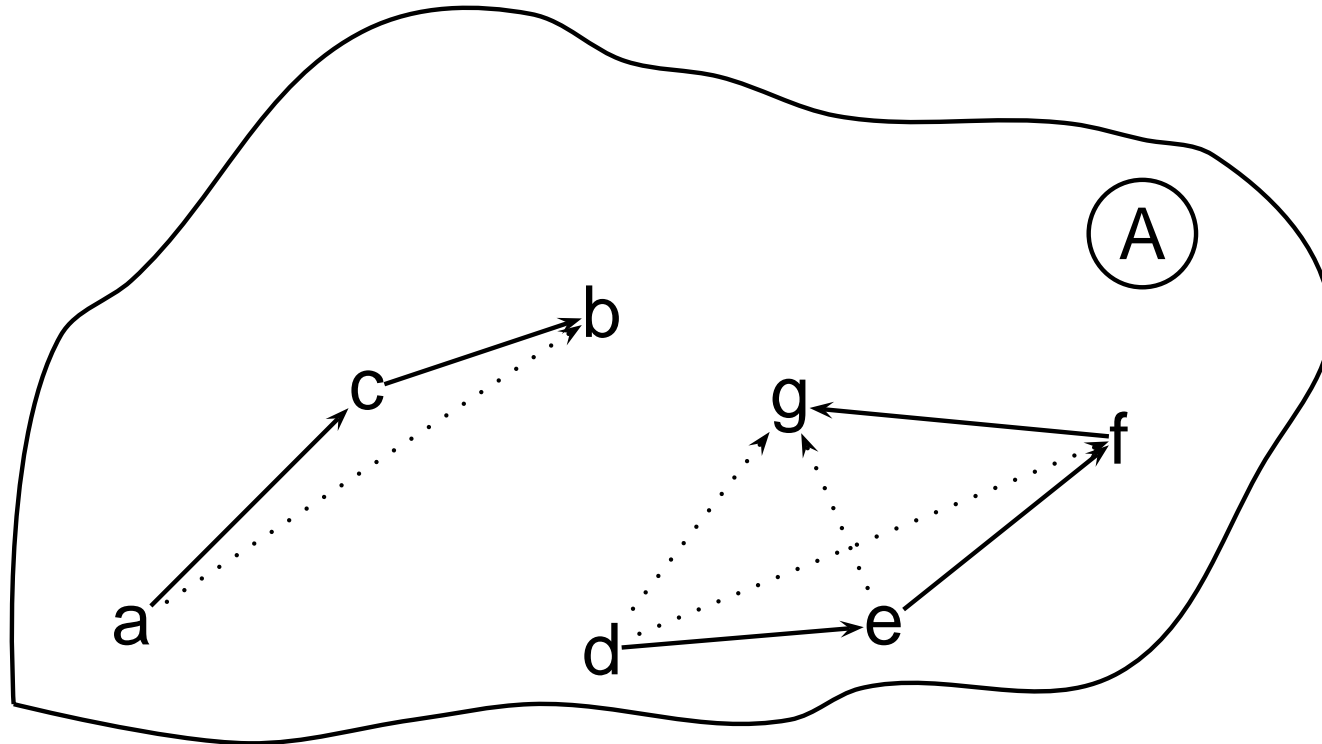
Solid links give a partial picture of a relation R from A to A . If R is transitive, then when the solid links exist, the dashed links must also exist.

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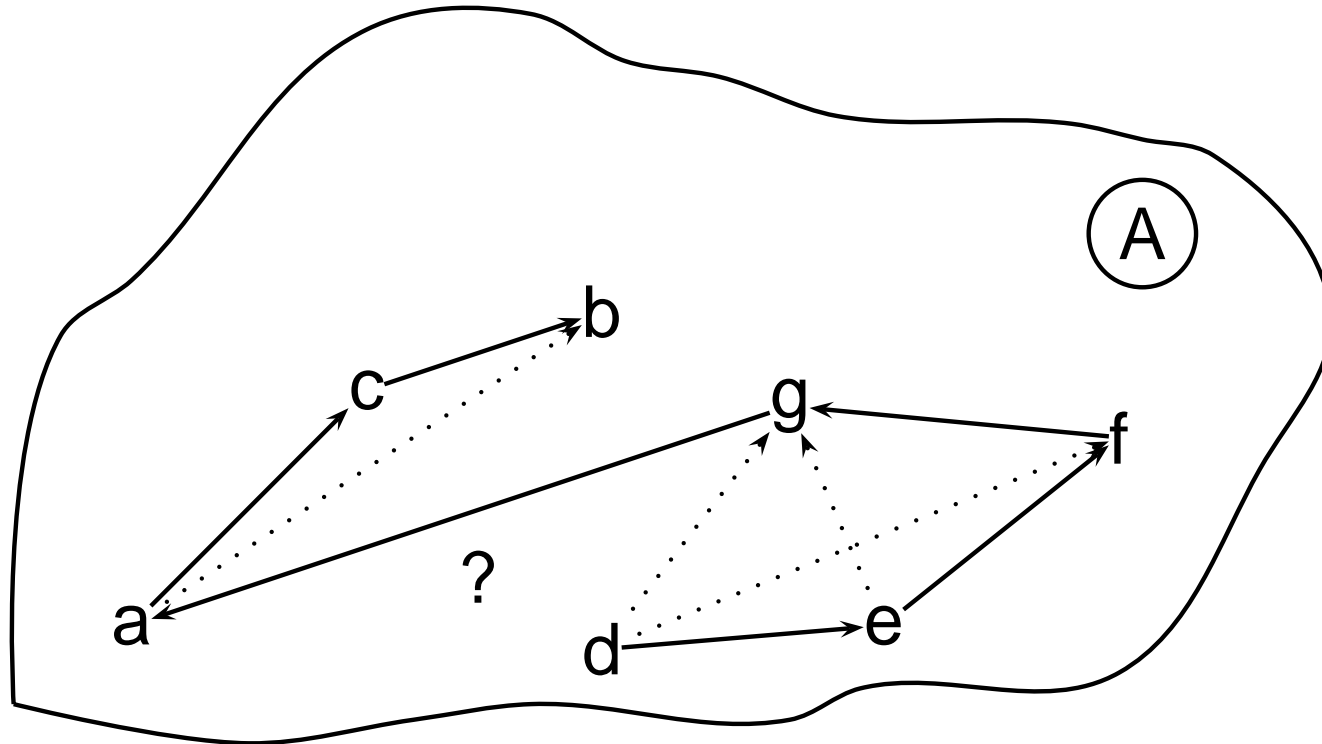
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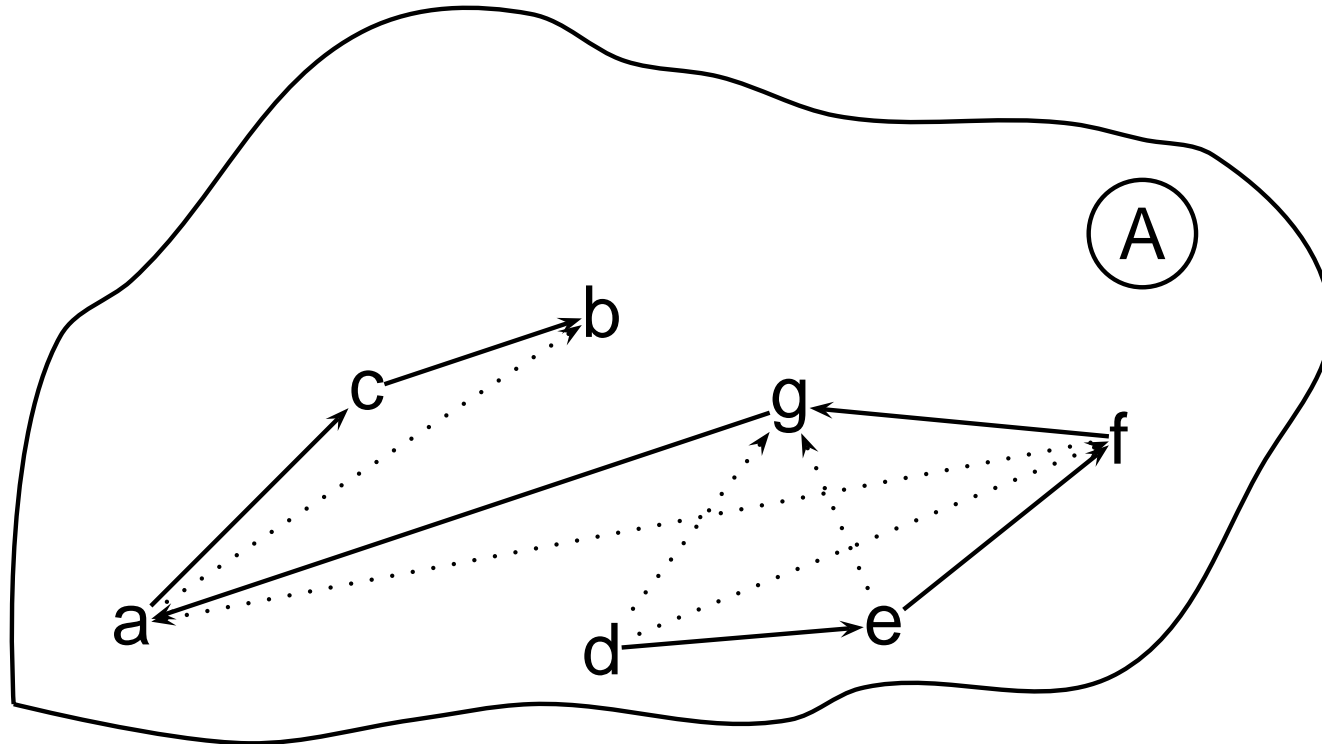
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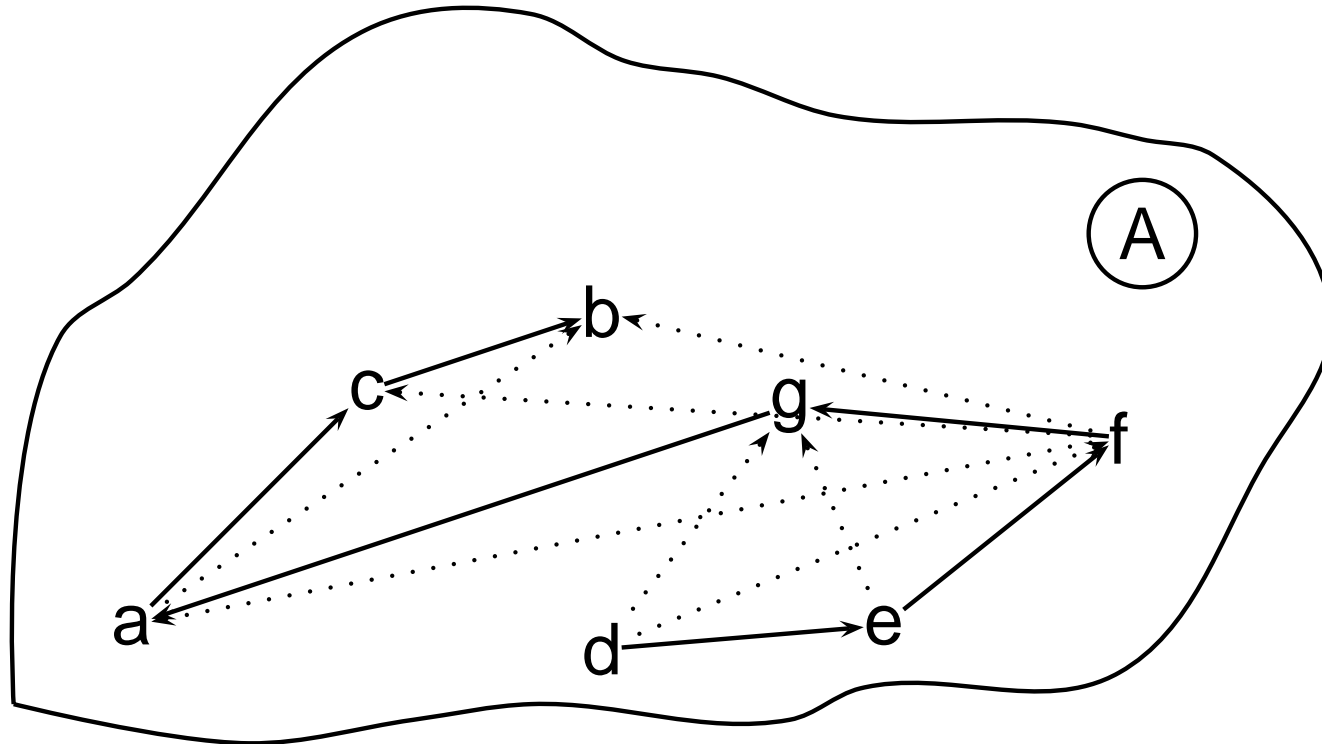
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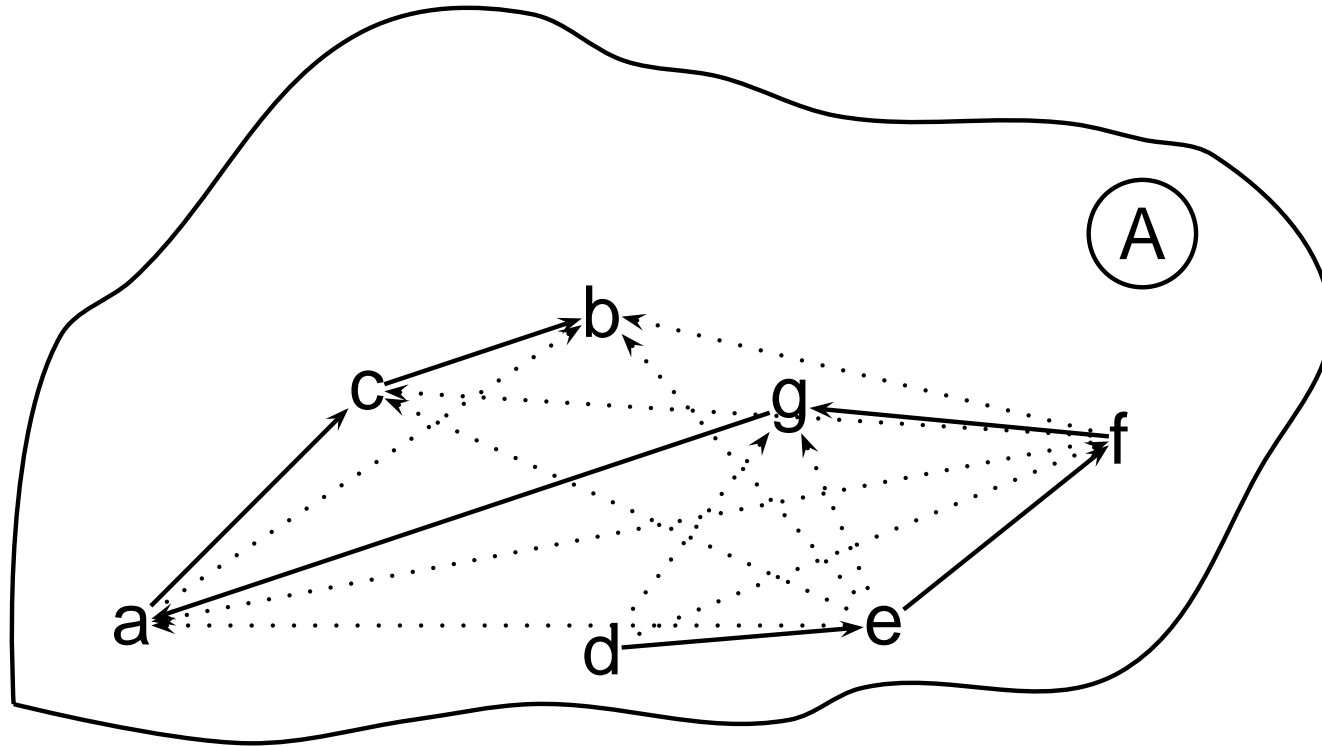
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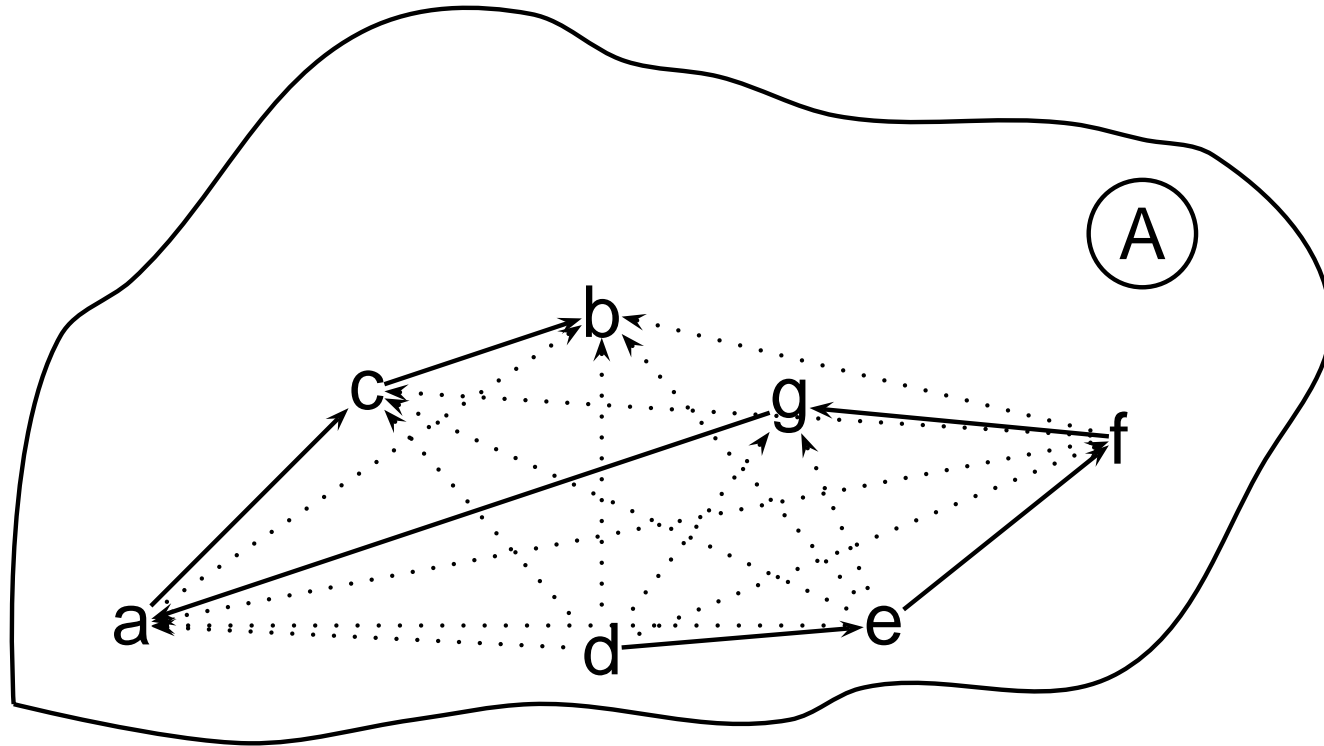
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Examples of transitivity

- The greater-than relation ($<$) on integers is transitive:

$$x < y$$

$$y < z$$

$$\Rightarrow$$

$$x < z$$

- The father-of relation is NOT transitive
Fred is Tom's father
Tom is Hank's father
 \nRightarrow
Fred is Hank's father
In fact. Fred is Hank's grandfather

More Transitivity examples

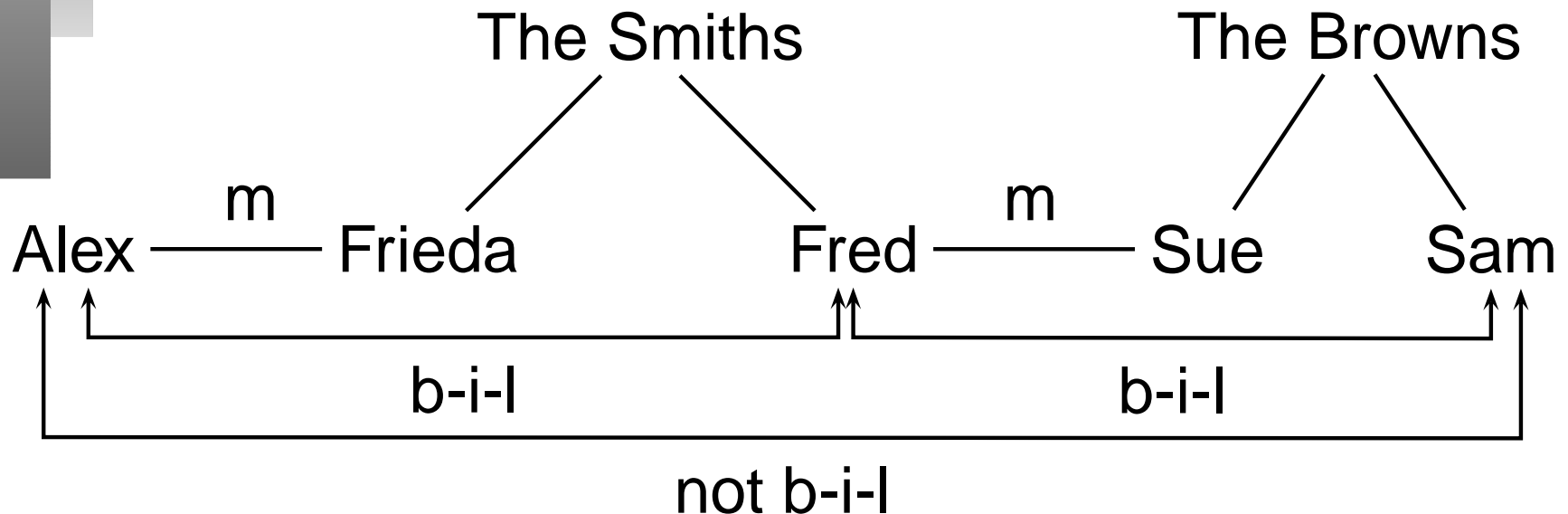
Example 14. *Transitivity of brother-of relation.*

If Fred is Sam's brother and Alex is is Fred's brother, then Alex is Sam's brother.

Example 15. *Non-Transitivity of brother-in-law relation.*

If Fred is Sam's brother-in-law (because Fred is married to Sam's sister Sue) and Alex is Fred's brother-in-law (say, because Alex is married to Fred's sister Frieda), then Alex is not Sam's brother-in-law.

Non-transitivity of brother-in-law relation



Non-transitivity of the *brother-in-law* relation. Frieda and Fred are syblings, as are Sue and Sam. Marriage links are labeled *m*; brother-in-law links are labeled *b-i-l*.

Definition 9. *Symmetry of a relation*
A relation R is symmetric if and only if

$$\forall \langle x, y \rangle \in R [\langle y, x \rangle \in R]$$

Symmetry examples

Example 16. *Symmetry of sybling relation*

If Fred is Sue's sybling then Sue is Fred's sybling.

Example 17. *Symmetry of cousin relation*

If Fred is Sue's cousin then Sue is Fred's cousin.

Example 18. *Non-symmetry of brother relation*

If Fred is Sue's brother and Sue is female it is not the case that Sue is Fred's brother.

Definition 10. *Asymmetry of a relation*
A relation R is asymmetric if and only if

$$\forall \langle x, y \rangle \in R [\langle y, x \rangle \notin R]$$

Example 19. *Asymmetry of father relation.*
If Fred is Bob's father it is definitely not the case that Bob is Fred's father.

Example 20. *Symmetry and the brother relation.*
If Fred is someone's brother they may or may not be his brother, depending on their sex.